

Analysis of Regular Sequences: Summatory Functions and Divide-and-Conquer Recurrences



Source and Acknowledgements

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Prelude 1: Find Min & Max via Divide & Conquer

- Partition into two sets of (almost) equal size;
- Find min and max in both parts individually & recursively;
- Compare minima, compare maxima.

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In other words:

$$M(2n) = 2M(n) + 2,$$

 $M(2n + 1) = M(n) + M(n + 1) + 2.$

Prelude 2: Binary Sum of Digits

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Clemens Heuberger

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q-Regular Sequences

Vector-valued sequence $\mathbf{v}\colon \mathbb{N}_0 \to \mathbb{C}^D$ with

$$v(qn+r)=A_rv(n)$$

for all $0 \le r < q$ and $n \ge 0$.

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Constants:

- $q \ge 2$, $D \ge 1$: integers;
- A_0, \ldots, A_{q-1} : $D \times D$ -matrices.

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First component of v: q-regular sequence (Allouche-Shallit 1992).

Theorem (Dumas 2013; H-Krenn 2020)

• x(n): q-regular sequence, first component of v(n)

$$\sum_{0 \le n < N} x(n) = \sum_{\substack{\lambda \in \sigma(C) \\ |\lambda| > R}} N^{\log_q \lambda} \sum_{0 \le k < m_C(\lambda)} \frac{(\log N)^k}{k!} \Phi_{\lambda k}(\{\log_q N\}) + O(N^{\log_q R}(\log N)^{\max\{m_C(\lambda) : |\lambda| = R\}})$$

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- $m_C(\lambda)$: size of the largest Jordan block of C associated with λ

$$\sum_{0 \le n < N} x(n) = \sum_{\substack{\lambda \in \sigma(C) \\ |\lambda| > R}} N^{\log_q \lambda} \sum_{0 \le k < m_C(\lambda)} \frac{(\log N)^k}{k!} \Phi_{\lambda k}(\{\log_q N\}) + O(N^{\log_q R}(\log N)^{\max\{m_C(\lambda) : |\lambda| = R\}})$$

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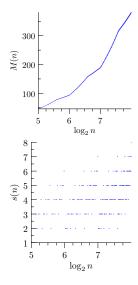
$$\Phi_{\lambda k}(u) = \sum_{\mu \in \mathbb{Z}} \varphi_{\lambda k \mu} \exp(2\mu \pi i u)$$

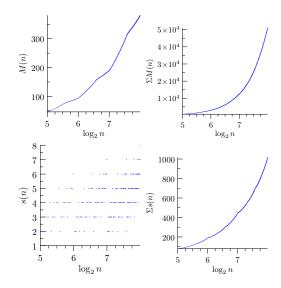
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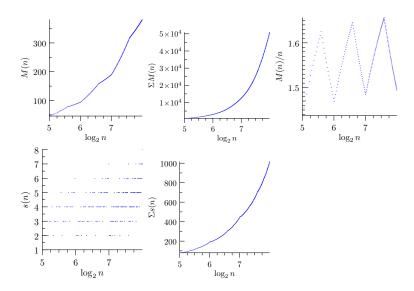
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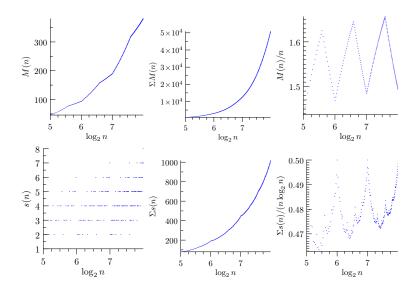
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 Fourier coefficients can be computed numerically (using a functional equation for the corresponding Dirichlet series)









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Spoilers

- Yes (almost always)
- For divide & conquer sequences

"good asymptotic expansion".

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Let x be a q-regular sequence with matrices $(A_r)_{0 \le r < q}$. Set $C := \sum_{0 \le r < q} A_r$. Assume that C has an eigenvalue $\neq 0$. Then there is a non-negative integer k such that $\Sigma^k x$ admits a "good asymptotic expansion".

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Repeating: asymptotic domination by D^kC (if C has an eigenvalue $\neq 0$).

Divide-and-Conquer Sequences

Divide-and-conquer sequence:

$$x(n) = \alpha x \left(\left\lfloor \frac{n}{2} \right\rfloor \right) + \beta x \left(\left\lceil \frac{n}{2} \right\rceil \right) + g(n)$$

for $n \ge 2$ (α , β given positive constants, g given function ("toll function"), x(1) given).

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Assume that there is an $\varepsilon > 0$ such that $g(n) = O(n^{\log_2(\alpha+\beta)-\varepsilon})$. Then

$$x(n) = n^{\log_2(\alpha+\beta)} \Phi(\{\log_2 n\}) + O(n^{\log_2(\alpha+\beta)-\varepsilon})$$

for $n \to \infty$ where Φ is a continuous, 1-periodic function.

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Question: Relation to "our" result?

$$x(n) = \alpha x \left(\left\lfloor \frac{n}{2} \right\rfloor \right) + \beta x \left(\left\lceil \frac{n}{2} \right\rceil \right) + g(n)$$

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$$\iff (\text{with } v(n) = (x(n), x(n+1))^{\top})$$

$$v(2n) = \begin{pmatrix} \alpha + \beta & 0 \\ \alpha & \beta \end{pmatrix} v(n) + \begin{pmatrix} g(2n) \\ g(2n+1) \end{pmatrix}$$

$$v(2n+1) = \begin{pmatrix} \alpha & \beta \\ 0 & \alpha + \beta \end{pmatrix} v(n) + \begin{pmatrix} g(2n+1) \\ g(2n+2) \end{pmatrix}$$

$$x(n) = \alpha x \left(\left\lfloor \frac{n}{2} \right\rfloor \right) + \beta x \left(\left\lceil \frac{n}{2} \right\rceil \right) + g(n)$$

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If g is regular, then x is regular.

Summatory Function of the Forward Difference

$$x(N) = x(0) + \sum_{0 \le n \le N} (x(n+1) - x(n))$$

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- Forward difference of regular sequence is regular
- Summatory function of regular sequence is regular
- ... but can we say something in general?

Set
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.

Set $\Delta x(n) := x(n+1) - x(n)$. Consider divide-and-conquer sequence

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 \Rightarrow

$$\Delta x(2n) = \beta \Delta x(n) + g(2n+1) - g(2n)$$

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$$\Delta x(2n+1) = \alpha \Delta x(n) + g(2n+2) - g(2n+1).$$

Dimension 1 (plus dimension of linear representation of g): particularly simple.

Theorem (H.–Krenn–Lechner 2024)

Let x be a divide-and-conquer sequence with polynomial toll function of degree $k \geq 1$. Then (for 1-periodic continuous functions Φ and Ψ and $n \rightarrow \infty$):

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• Case 1a. If $\alpha + \beta > 2^k$ and $2^k > \max\{\alpha, \beta\}$, then

$$\begin{split} x(\mathbf{n}) &= \mathbf{n}^{\log_2(\alpha+\beta)} \Phi(\{\log_2 \mathbf{n}\}) + \mathbf{n}^k \Psi(\{\log_2 \mathbf{n}\}) \\ &+ O(\mathbf{n}^{\log_2 \max\{\alpha,\beta\}}). \end{split}$$

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• Case 1b. If $\alpha + \beta > 2^k$ and $\max\{\alpha, \beta\} \ge 2^k$, then

$$\begin{split} x(\textbf{n}) &= \textbf{n}^{\log_2(\alpha+\beta)} \Phi(\{\log_2 \textbf{n}\}) \\ &+ O(\textbf{n}^{\log_2 \max\{\alpha,\beta\}} (\log \textbf{n})^{[\max\{\alpha,\beta\}=2^k]}). \end{split}$$

Theorem (H.-Krenn-Lechner 2024)

Let x be a divide-and-conquer sequence with polynomial toll function of degree $k \geq 1$. Then (for 1-periodic continuous functions Φ and Ψ and $n \rightarrow \infty$):

• Case 2. If $\alpha + \beta = 2^k$, then

$$x(n) = n^{k}(\log n)\Phi(\{\log_{2} n\}) + n^{k}\Psi(\{\log_{2} n\}) + O(n^{\log_{2} \max\{\alpha,\beta\} + [\alpha = \beta]\varepsilon})$$

for any $\varepsilon > 0$.

Theorem (H.-Krenn-Lechner 2024)

Let x be a divide-and-conquer sequence with polynomial toll function of degree $k \geq 1$. Then (for 1-periodic continuous functions Φ and Ψ and $n \rightarrow \infty$):

• Case 3. If $2^k > \alpha + \beta > 2^{k-1}$, then

$$\begin{split} x(\textbf{\textit{n}}) &= \textbf{\textit{n}}^k \Phi(\{\log_2 \textbf{\textit{n}}\}) + \textbf{\textit{n}}^{\log_2(\alpha+\beta)} \Psi(\{\log_2 \textbf{\textit{n}}\}) \\ &+ O\Big(\textbf{\textit{n}}^{\log_2 \max\{\alpha,\beta,2^{k-1}\} + [\max\{\alpha,\beta\} = 2^{k-1}]\varepsilon} \\ &\times \big(\log \textbf{\textit{n}}\big)^{[\max\{\alpha,\beta\} < 2^{k-1}]}\Big) \end{split}$$

for any $\varepsilon > 0$.

Theorem (H.-Krenn-Lechner 2024)

Let x be a divide-and-conquer sequence with polynomial toll function of degree $k \geq 1$. Then (for 1-periodic continuous functions Φ and Ψ and $n \rightarrow \infty$):

• Case 4. If $2^{k-1} \ge \alpha + \beta$, then

$$x(\mathbf{n}) = \mathbf{n}^k \Phi(\{\log_2 \mathbf{n}\}) + O(\mathbf{n}^{k-1}(\log \mathbf{n})^E),$$

where

$$E := 1 + [\alpha + \beta = 2^{k-1}]([k \ge 2 \text{ and } c_{k-1} \ne 0] + [k = 1 \text{ and } d_0 + d_1 \ne 0])$$

with

$$d_0 := (1 - \beta)x(1) - g(1) + g(0), \quad d_1 := g(1) - (1 - \beta)x(1).$$

Min-Max: Fourier Coefficients

Computing Fourier coefficients using the general result . . .

