

# The Giant Component and 2-Core in Sparse Random Outerplanar Graphs

**Michael Missethan**

(Joint work with Mihyun Kang)

31st International Conference on Probabilistic, Combinatorial and Asymptotic Methods for the Analysis of Algorithms 2020



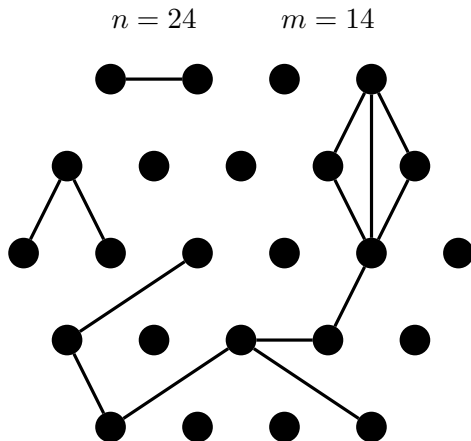
- Random planar graphs
- Emergence of the giant component
  - Random planar graphs
  - Random outerplanar graphs
- Core–kernel approach for planar graphs
- Core approach for outerplanar graphs

# Random planar graph $P(n, m)$

- Pick a planar graph with  $n$  vertices and  $m$  edges (uniformly at random).

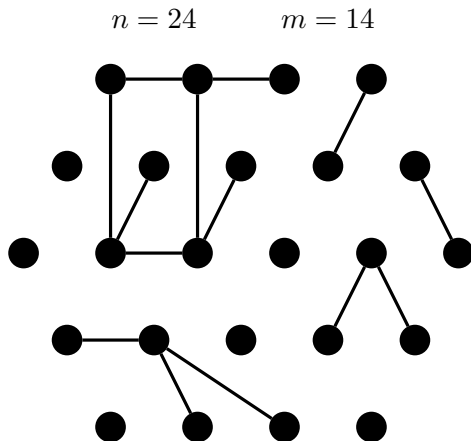
# Random planar graph $P(n, m)$

- Pick a planar graph with  $n$  vertices and  $m$  edges (uniformly at random).



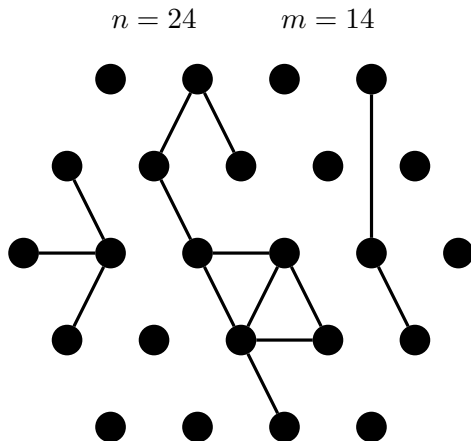
# Random planar graph $P(n, m)$

- Pick a planar graph with  $n$  vertices and  $m$  edges (uniformly at random).



# Random planar graph $P(n, m)$

- Pick a planar graph with  $n$  vertices and  $m$  edges (uniformly at random).



# Random planar graph $P(n, m)$

- $m = m(n)$  (e.g.  $n/2, n, \dots$ )
- $n \rightarrow \infty$

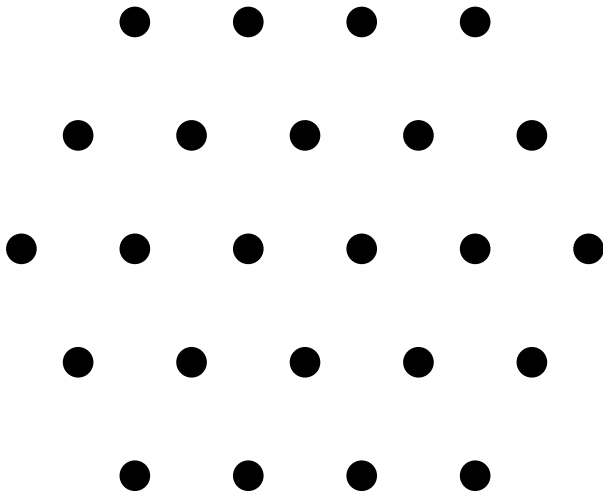
# Random planar graph $P(n, m)$

- $m = m(n)$  (e.g.  $n/2, n, \dots$ )
- $n \rightarrow \infty$

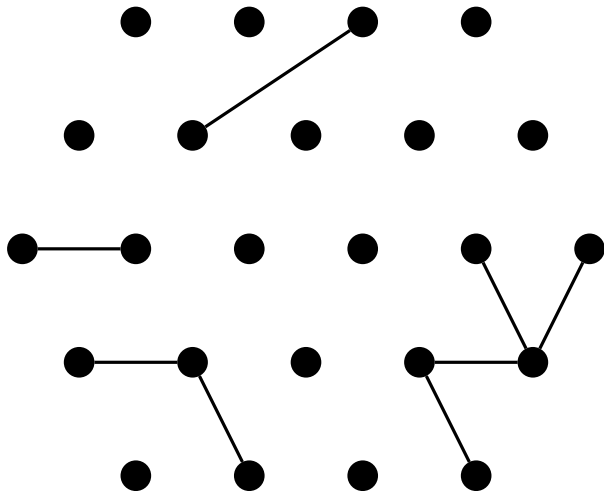
- How does the component structure change when  $m$  varies?



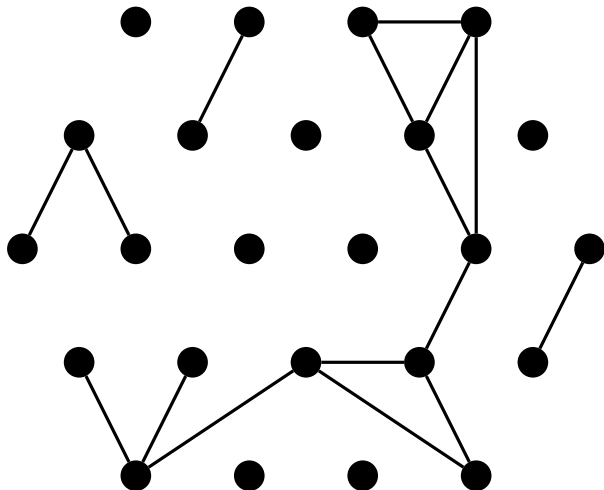
# Random planar graph $P(n, m)$



# Random planar graph $P(n, m)$



# Random planar graph $P(n, m)$



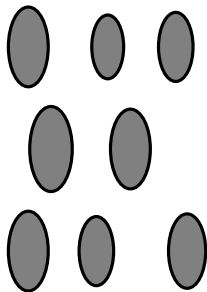
# Random planar graph $P(n, m)$

## Theorem

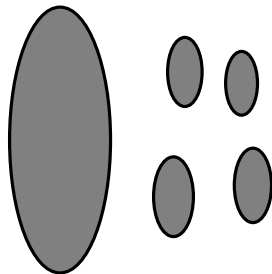
[Kang, Łuczak, 2012]

The giant component emerges at

$$m = n/2 + O(n^{2/3}).$$



$$m \leq n/2 + O(n^{2/3})$$

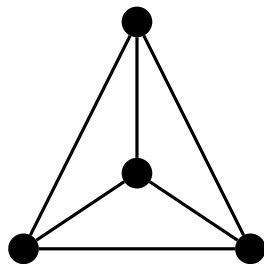
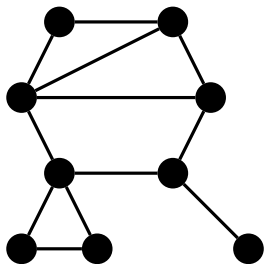


$$m \geq n/2 + \omega(n^{2/3})$$

# Random outerplanar graphs

## Outerplanar graph

has a planar drawing, such that all vertices belong to the outer face.



## Theorem

[Kang-M. 2020]

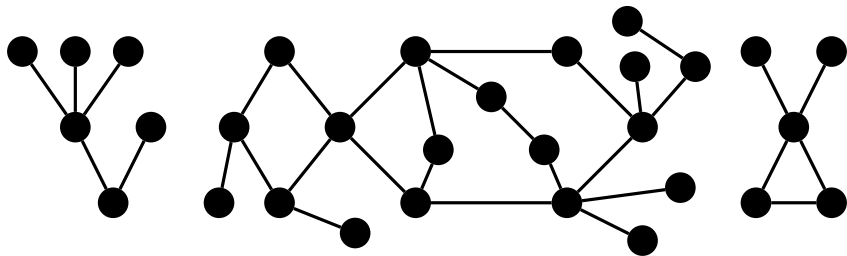
Outerplanar graphs feature a similar phase transition as planar graphs, i.e. the giant component emerges at

$$m = n/2 + O\left(n^{2/3}\right).$$

# Core–kernel approach for planar graphs

## Decomposition

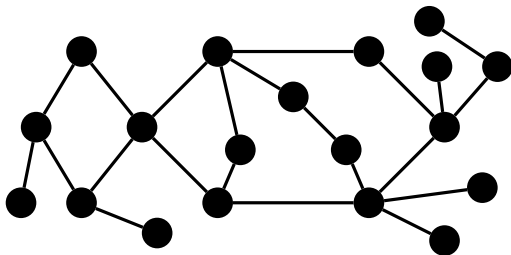
- *Complex part*: union of components with at least two cycles



# Core–kernel approach for planar graphs

## Decomposition

- *Complex part*: union of components with at least two cycles

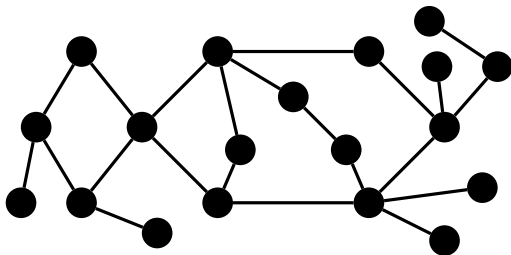




# Core–kernel approach for planar graphs

## Decomposition

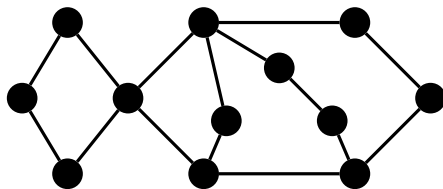
- *Complex part*: union of components with at least two cycles
- *Core*: maximal subgraph of minimum degree at least two



# Core–kernel approach for planar graphs

## Decomposition

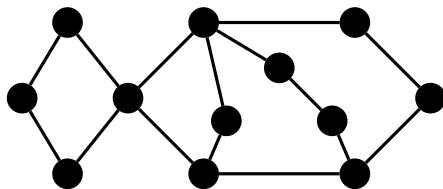
- *Complex part*: union of components with at least two cycles
- *Core*: maximal subgraph of minimum degree at least two



# Core–kernel approach for planar graphs

## Decomposition

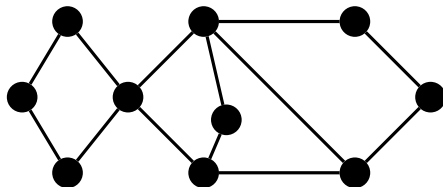
- *Complex part*: union of components with at least two cycles
- *Core*: maximal subgraph of minimum degree at least two
- *Kernel*: replace paths consisting of vertices of degree two by edges



# Core–kernel approach for planar graphs

## Decomposition

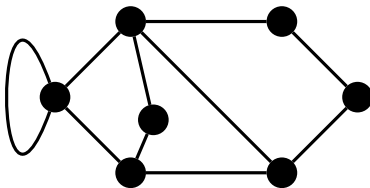
- *Complex part*: union of components with at least two cycles
- *Core*: maximal subgraph of minimum degree at least two
- *Kernel*: replace paths consisting of vertices of degree two by edges



# Core–kernel approach for planar graphs

## Decomposition

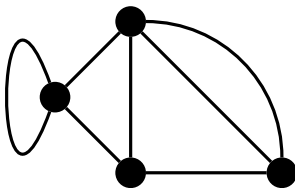
- *Complex part*: union of components with at least two cycles
- *Core*: maximal subgraph of minimum degree at least two
- *Kernel*: replace paths consisting of vertices of degree two by edges



# Core–kernel approach for planar graphs

## Decomposition

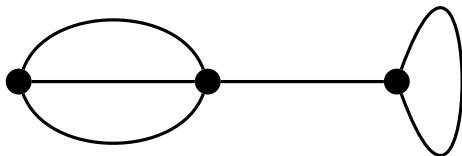
- *Complex part*: union of components with at least two cycles
- *Core*: maximal subgraph of minimum degree at least two
- *Kernel*: replace paths consisting of vertices of degree two by edges



# Core–kernel approach for planar graphs

## Construction

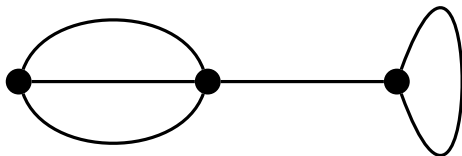
- Start with a kernel



# Core–kernel approach for planar graphs

## Construction

- Start with a kernel
- Subdivide edges to obtain core

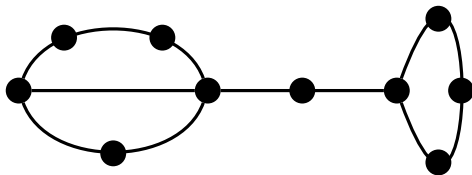




# Core–kernel approach for planar graphs

## Construction

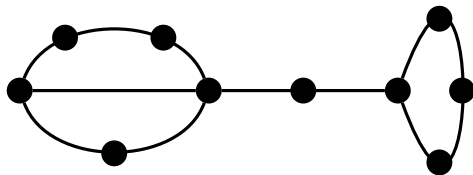
- Start with a kernel
- Subdivide edges to obtain core



# Core–kernel approach for planar graphs

## Construction

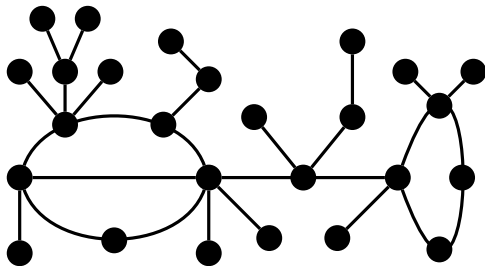
- Start with a kernel
- Subdivide edges to obtain core
- Attach rooted trees to get complex part



# Core–kernel approach for planar graphs

## Construction

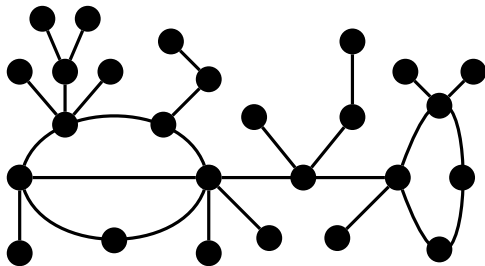
- Start with a kernel
- Subdivide edges to obtain core
- Attach rooted trees to get complex part



# Core–kernel approach for planar graphs

## Construction

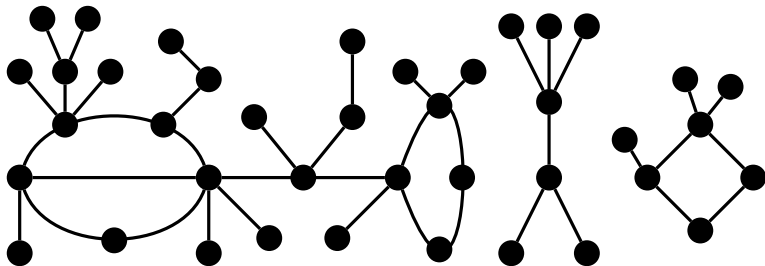
- Start with a kernel
- Subdivide edges to obtain core
- Attach rooted trees to get complex part
- Add trees and unicyclic components



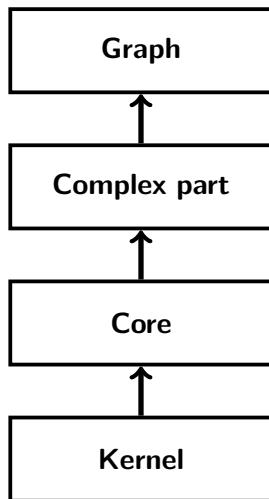
# Core–kernel approach for planar graphs

## Construction

- Start with a kernel
- Subdivide edges to obtain core
- Attach rooted trees to get complex part
- Add trees and unicyclic components



# Core–kernel approach for planar graphs

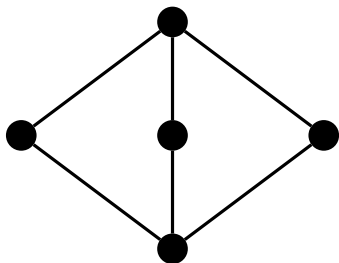


# Core–kernel approach for planar graphs

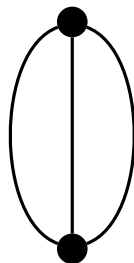
- Key fact: a graph is planar  $\iff$  its kernel is planar

# Core–kernel approach for planar graphs

- Key fact: a graph is planar  $\iff$  its kernel is planar
- BUT: A non-outerplanar graph can have an outerplanar kernel.



$K_{2,3}$



kernel of  $K_{2,3}$



# Core approach for outerplanar graphs

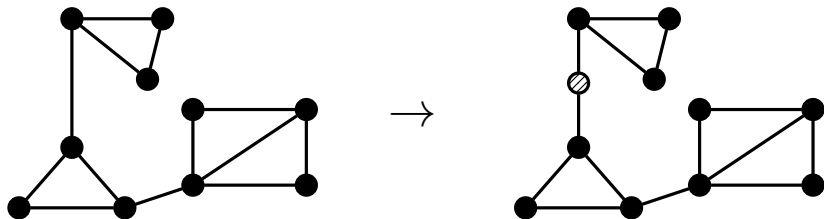
- A graph is outerplanar  $\iff$  its core is outerplanar

# Core approach for outerplanar graphs

- A graph is outerplanar  $\iff$  its core is outerplanar
- Determine  $C(n, m) = \# \text{cores with } n \text{ vertices and } m \text{ edges}$

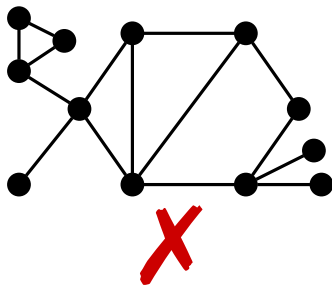
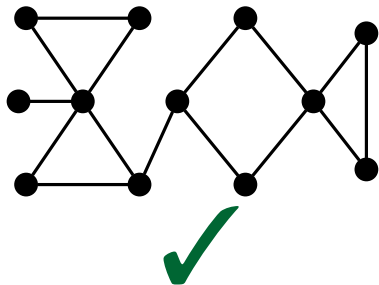
# Core approach for outerplanar graphs

- A graph is outerplanar  $\iff$  its core is outerplanar
- Determine  $C(n, m) = \#$ cores with  $n$  vertices and  $m$  edges
  - Estimate  $\frac{C(n+1, m+1)}{C(n, m)}$



# Core approach for outerplanar graphs

- A graph is outerplanar  $\iff$  its core is outerplanar
- Determine  $C(n, m) = \#$ cores with  $n$  vertices and  $m$  edges
  - Estimate  $\frac{C(n+1, m+1)}{C(n, m)}$
  - Lower bound: Cactus graphs



# Core approach for outerplanar graphs

- A graph is outerplanar  $\iff$  its core is outerplanar
- Determine  $C(n, m) = \# \text{cores with } n \text{ vertices and } m \text{ edges}$ 
  - Estimate  $\frac{C(n+1, m+1)}{C(n, m)}$
  - Lower bound: Cactus graphs
  - Upper bound: Planar graphs

- Emergence of giant component at  $m = n/2 + O(n^{2/3})$

- Emergence of giant component at  $m = n/2 + O(n^{2/3})$
- Proof idea for planar graphs: Core–kernel approach
  - Key fact: a graph is planar  $\iff$  its kernel is planar

- Emergence of giant component at  $m = n/2 + O(n^{2/3})$
- Proof idea for planar graphs: Core–kernel approach
  - Key fact: a graph is planar  $\iff$  its kernel is planar
- Proof idea for outerplanar graphs: Core approach
  - Direct analysis of the core (without using the kernel)